How to split Recursive Automata

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Available Information

- a set of positive examples
- the target class

First possible strategy: learning by generalization

- build a least general grammar generating the examples
- apply a generalization operator until it belongs to the target class

Second possible strategy: learning by specialization

- the initial hypothesis space is the whole target class
- use the examples to constrain this space until it is reduced to one grammar

Overview of known results

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|--------------------|-----------------------|---------------------|
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| representation | finite state automata | Categorial Grammars |

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| strategy | (Angluin 81) | (Kanazawa 96, 98) |

The links between them: in (Tellier 05, 06)

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The links between them: this paper!

Outline

- 1. Introduction
- 2. Categorial Grammars and Recursive Automata
- 3. Learning by specialization in both representations
- 4. Learning from Typed Examples: a new interpretation
- 5. Conclusion

Categorial Grammars and Recursive Automata

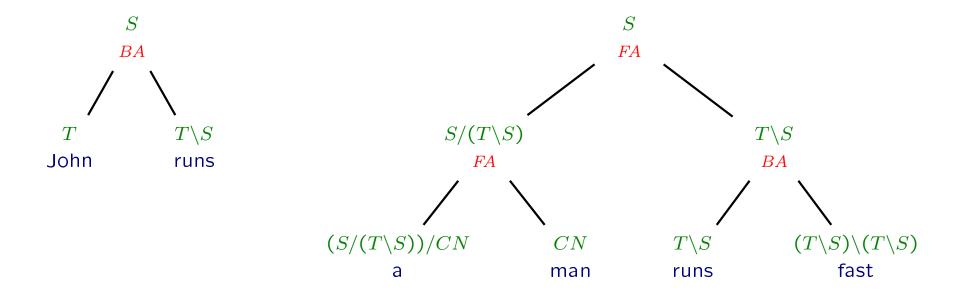
Definition of a AB-Categorial Grammar

- let Σ be a finite vocabulary
- let $\mathcal B$ be an enumerable set of basic categories, among which is the axiom $S \in \mathcal B$
- the set of categories $Cat(\mathcal{B})$ is the smallest set such that :
 - $-\mathcal{B}\subset Cat(\mathcal{B})$
 - $\forall A, B \in Cat(\mathcal{B}) : A/B \in Cat(\mathcal{B}) \text{ and } B \setminus A \in Cat(\mathcal{B})$
- a Categorial Grammar G is a finite relation over $\Sigma \times Cat(\mathcal{B})$
- Syntactic rules are expressed by two schemes : $\forall A, B \in Cat(\mathcal{B})$
 - Forward Application $FA: A/B \ B \longrightarrow A$
 - Backward Application $BA: B B \setminus A \longrightarrow A$
- $-\ L(G)$: set of strings corresponding to a sequence of categories which reduces to S

Categorial Grammars and Recursive Automata

Definition of a AB-Categorial Grammar

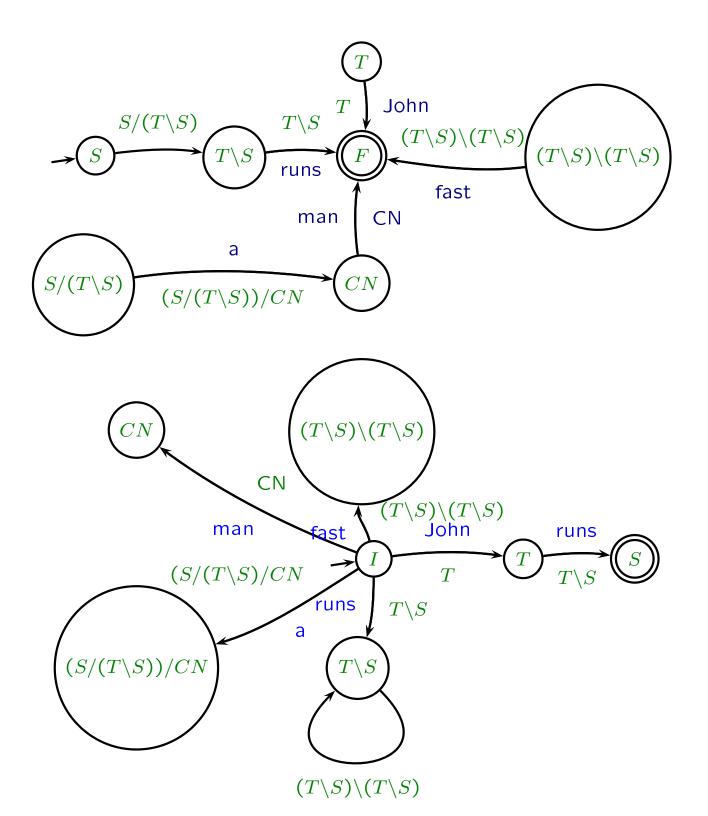
- Let $\mathcal{B} = \{S, T, CN\}$ where T stands for "term" and CN for "common noun"
- $-\Sigma = \{John, runs, a, man, fast\}$
- $-G = \{\langle \mathsf{John}, T \rangle, \langle \mathsf{runs}, T \backslash S \rangle, \langle \mathsf{a}, (S/(T \backslash S))/CN \rangle, \langle \mathsf{man}, CN \rangle \\ \langle \mathsf{fast}, (T \backslash S) \backslash (T \backslash S) \rangle \}$



Categorial Grammars and Recursive Automata

Definition of Recursive Automata (Tellier06)

- A RA is like a Finite State Automaton except that transitions can be labelled by a state
- Using a transition labelled by a state Q means producing $w \in L(Q)$
- There are two distinct kinds of RA:
 - the RA_{FA} -kind where the language L(Q) of a state Q is the set of strings from Q to the final state
 - Every unidirect. FA CG is strongly equivalent with a RA_{FA}
 - the RA_{BA} -kind where the language L(Q) of a state Q is the set of strings from the initial state to Q
 - Every unidirect. BA CG is strongly equivalent with a RA_{BA}
- Every CG is equivalent with a pair $MRA = \langle RA_{FA}, RA_{FA} \rangle$



Outline

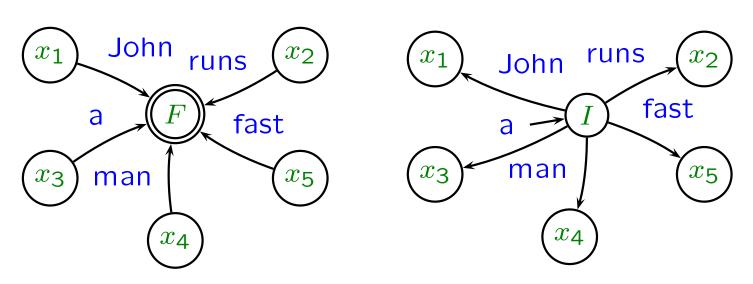
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Inference of rigid CGs from strings (Moreau 04)

- Target Class: rigid Categorial Grammars, i.e. at most one category for each word
- Input: a set of sentences
- Learning Algorithm :
 - 1. associate a distinct unique variable with each word
 - 2. for each sentence do
 - try to parse the sentence (CYK-like algorithm)
 - induce constraints on the variables
- Output: (disjunctions of) set(s) of constraints, each set
 corresponding with a (set of) rigid grammar(s)

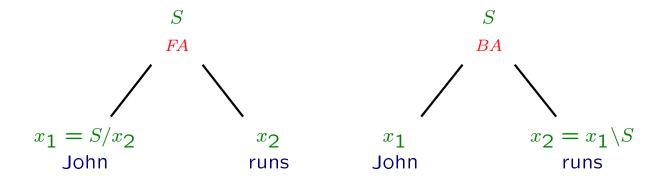
Inference of rigid CGs from strings (Moreau 04): example

- input data : The set $D = \{John runs, a man runs fast\}$
- associate a distinct unique variable with each word : $\mathcal{A} = \{ \langle \mathsf{John}, \ x_1 \rangle, \ \langle \mathsf{runs}, \ x_2 \rangle, \ \langle \mathsf{a}, \ x_3 \rangle, \ \langle \mathsf{man}, \ x_4 \rangle, \ \langle \mathsf{fast}, \ x_5 \rangle \}$
- for every unidirectional CG G, there exists a substitution transforming ${\mathcal A}$ into G
- \mathcal{A} specifies the set of every unidirectional CGs
- \mathcal{A} can also be represented by a $MRA = \langle RA_{FA}, RA_{BA} \rangle$:



Inference of rigid CGs from strings (Moreau 04): example

— the only two possible ways to parse "John runs" :



- to parse "a man runs fast" :
 - theoretically : $5 * 2^3 = 40$ distinct possible ways
 - but some couples of constraints are not compatible with the class of rigid grammars
- main problem with this algo: combinatorial explosion
- to limit it: initial knowledge in the form of known assignments

Effects of constraints on a $MRA = \langle RA_{FA}, RA_{BA} \rangle$

- constraints inferred are of the form :
 - $x_k=x_l$ with x_k and x_l already exist : state and/or transition merges in both the RA_{FA} and the RA_{BA}
 - or $x_k = X_m/X_n$ (resp. $x_k = X_m \setminus X_n$) with $X_m, X_n \in Cat(\mathcal{B})$
- the effect of $x_k = X_m/X_n$ (resp. $x_k = X_m \setminus X_n$) in a MRA :
 - $-X_m/X_n$ (resp. $x_k=X_m\backslash X_n$) replaces x_k everywhere in the MRA
 - every subcategory of X_m and X_n (including themselves) becomes a new state in both the RA_{FA} and the RA_{BA} , linked to F (resp. from I) by a its name
 - in the RA_{FA} (resp. the RA_{BA}), a new transition labelled by X_m/X_n (resp. $X_m\backslash X_n$) links the states X_m and X_n
 - the states of the same name are merged
- So: a combination of state splits and state merges
- better founded than the state splits in (Fredouille 00)

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Basic ideas (Dudau, Tellier & Tommasi 01)

- cognitive hypothesis : lexical semantics is learned before syntax
- formalization : words are given with their (Montague's) semantic type
- Types derive from categories by a homomorphism
- Classical example : $h(T)=e,\ h(S)=t,\ h(CN)=\langle e,t\rangle$ and $h(A/B)=h(B\backslash A)=\langle h(B),h(A)\rangle$
- input data : typed sentences are of the form

| John | runs | а | man | runs | fast |
|------|----------------------|------------------------------------------------------------------------------|-----------------------|-----------------------|-----------------------------------------------------------|
| e | $\langle e,t angle$ | $\langle\langle e, t \rangle, \langle\langle e, t \rangle, t \rangle\rangle$ | $\langle e,t \rangle$ | $\langle e,t \rangle$ | $\langle\langle e,t \rangle, \langle e,t \rangle \rangle$ |

Target Class

- The set of CGs such that every distinct category assigned to the same word gives a distinct type
- $\forall \langle v, C_1 \rangle, \langle v, C_2 \rangle \in G, C_1 \neq C_2 \Longrightarrow h(C_1) \neq h(C_2)$
- Theorem (Dudau, Tellier & Tommasi 03): for every
 CF-language, there exists a grammar G generating it and a morphism h satisfying this condition

General algorithm (Dudau, Tellier & Tommasi 01)

- 1. initial set of assignments: introduce variables to represent the class
- 2. for each sentence
 - try to parse the sentence (CYK-like)
 - induce constraints on the variables
- 3. Output: (disjunctions) of set(s) of contraint(s), each being represented by a least general grammar

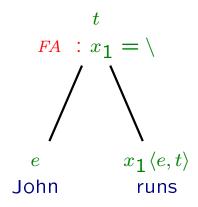
Example of pre-treatment

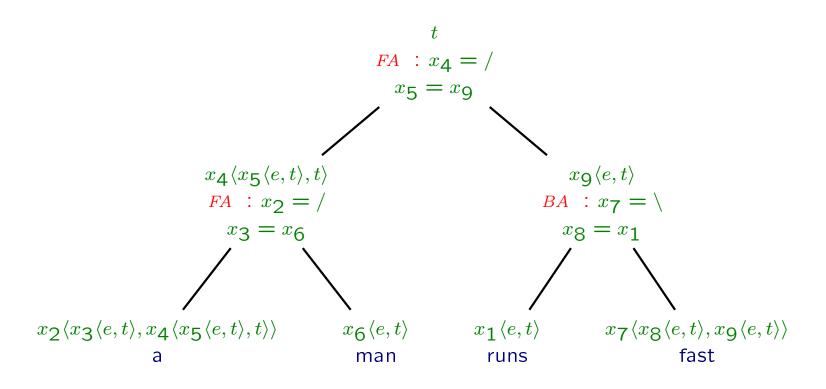
- introduce a distinct variable whose possible values are / or \backslash in front of every subtype
- in our example, the result is of the form :

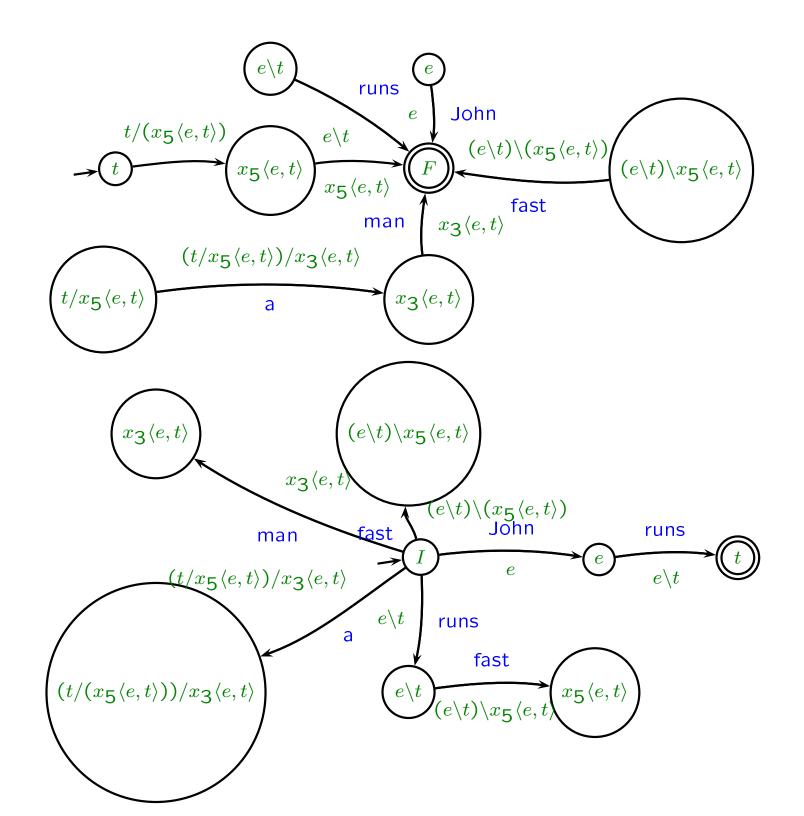
John runs
$$e x_1 \langle e, t \rangle$$

| а | man | runs | fast |
|--------------------------------------------------------------------------------------|-------------------------|-------------------------|------------------------------------------------------------------|
| $x_2\langle x_3\langle e,t\rangle, x_4\langle x_5\langle e,t\rangle,t\rangle\rangle$ | $x_6\langle e,t\rangle$ | $x_1\langle e,t\rangle$ | $x_7\langle x_8\langle e,t\rangle, x_9\langle e,t\rangle\rangle$ |

Infering constraints by parsing







Sum-up

- mix of state splits and state merges
- Types contain in themselves where splits are possible
- not every (complex) state can be merged: states are typed in the sense of (Coste & alii 2004)
- the use of types reduces the combinatorial explosion of possible splits
- types helph to converge to the correct solution quicker

Sum-up

| vocabulary | Moreau's initial | target category | pre-treated type |
|------------|------------------|--------------------------|--------------------------------------------------------------------------------------|
| | assigment | | |
| а | x_1 | $(S/(T\backslash S))/CN$ | $x_2\langle x_3\langle e,t\rangle, x_4\langle x_5\langle e,t\rangle,t\rangle\rangle$ |
| man | x_2 | CN | $x_6\langle e, t \rangle$ |
| runs | x_3 | $T \backslash S$ | $x_1\langle e,t\rangle$ |

- there exists a substitution, thus a homomorphism between
 Moreau's assignments and categories
- there exists a homomorphism between categories and types
 (Principle of compositionality)
- the starting point is either a lower bound or an upper bound
- the "good substitution" is well constrained by types

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Conclusion

Main contributions

- we mainly propose a new perspective on already known algorithms
- the correspondence between Categorial Grammars and recursive automata is fruitful
- MRA can represent sets of grammars corresponding to search spaces
- specialization strategies require additional knowledge (like semantic types)
- natural language is probably learnt by specialization by children
- specialization techniques deserve further investigation (better for incrementality...)