

Relevant Representations for the Inference of Rational Stochastic Tree Languages

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Outline

- 1 The Basic Problem
- 2 A Canonical Linear Representation for Rational Tree Series
- 3 Contributions
 - Normalization of the Model as a Generative Model
 - Strongly Consistent Model
 - Unranked Trees

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Trees

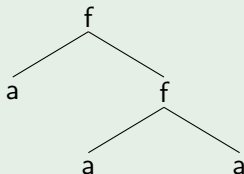
$\mathcal{F} = \mathcal{F}_0 \cup \mathcal{F}_1 \cup \dots \cup \mathcal{F}_p$: a ranked alphabet

\mathcal{F}_m : function symbols of *arity* m .

$T(\mathcal{F})$: all the *trees* constructed from \mathcal{F} .

Example:

$\mathcal{F} = \{f(\cdot, \cdot), a\}$; $f(a, f(a, a)) \in T(\mathcal{F})$.



Stochastic Tree Languages

Stochastic tree language: Probability distribution over $T(\mathcal{F})$

$$p : T(\mathcal{F}) \rightarrow \mathbb{R}$$

- for any $t \in T(\mathcal{F})$, $0 \leq p(t) \leq 1$ and
- $\sum_{t \in T(\mathcal{F})} p(t) = 1$.

Formal power tree series over $T(\mathcal{F})$

$$r : T(\mathcal{F}) \rightarrow \mathbb{R}.$$

Notation: $\mathbb{R}\langle\langle T(\mathcal{F}) \rangle\rangle$ (vector space).

A Basic Problem in Probabilistic Grammatical Inference

The Problem

Data $t_1, \dots, t_n \in T(\mathcal{F})$ independently drawn according to a fixed unknown stochastic tree language p .

Goal Infer an estimate of p in some class of probabilistic models.

Probabilistic models

- Probabilistic tree automata
- Linear representations of rational tree series

Probabilistic Tree Automata

A distribution over $T(\mathcal{F})$ according to a PA with one state

$$\mathcal{A}_\alpha : \Delta_\alpha = \{q \xrightarrow{\alpha} a, q \xrightarrow{1-\alpha} f(q, q)\}, \quad \tau(q) = 1, \quad 0 \leq \alpha \leq 1$$

$$p_\alpha(f(a, f(a, a))) = \alpha^3(1 - \alpha)^2$$

Less simple than in the word case

- p_α is a stochastic language iff $\alpha \geq 1/2$.
- Is it decidable whether a PA defines a stochastic language?
- The average tree size: $1/(2\alpha - 1)$. Unbounded if $\alpha = 1/2$.
- It is polynomially decidable whether a PA defines a stochastic language with bounded average size.

Linear Representations of Rational Tree Languages

A series $r \in \mathbb{R}\langle\langle T(\mathcal{F}) \rangle\rangle$ is **rational** iff there exists a triple (V, μ, λ) :

- V is a finite dimensional vector space over \mathbb{R} ,
- μ maps any $f \in \mathcal{F}_p$ to a p -linear mapping $\mu(f) \in \mathcal{L}(V^p; V)$,
- λ is a linear form $V \rightarrow \mathbb{R}$,
- $r(t) = \lambda\mu(t)$, where $\mu(f(t_1, \dots, t_p)) = \mu(f)(\mu(t_1), \dots, \mu(t_p))$.

Example

- $V = \mathbb{R}$ and let $e_1 \neq 0$ a basis of \mathbb{R} ,
- $\mu(a) = \alpha e_1$, $\mu(f)(e_1, e_1) = (1 - \alpha)e_1$,
- $\lambda(e_1) = 1$.

$$\lambda\mu(f(a, f(a, a))) = \alpha^3(1 - \alpha)^2$$

Rational Stochastic Tree Languages

Stochastic languages

A *rational stochastic tree language (RSTL)* is a stochastic language that has a linear representation.

- Every stochastic language computed by a **probabilistic automaton** is rational.
- Some RSTL **cannot be computed** by a probabilistic automaton.
- It is **undecidable** whether a linear representation represents a stochastic language.
- A RSTL can be equivalently represented by a **weighted tree automaton**, minimal in the number of states (vector space).

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Word Languages: The Notion of Residual Languages

Languages: $L \subseteq \Sigma^*, u \in \Sigma^*$

$$u^{-1}L = \{v \in \Sigma^* \mid uv \in L\}$$

Series: $r \in \mathbb{R}\langle\langle T(\mathcal{F}) \rangle\rangle, u \in \Sigma^*$

$$\dot{u}r(v) = r(uv)$$

Residual language is a **key notion for inference** because:

- residual languages are **intrinsic components**
- they are **observable** on samples
- they yield **canonical representations**.

Contexts

$\$$: a zero arity function symbol not in \mathcal{F}_0 .

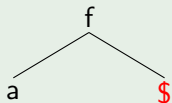
A *context* is an element of $T(\mathcal{F} \cup \{\$\})$ s.t. $\$$ appears exactly once.

$C(\mathcal{F})$: all contexts over \mathcal{F} .

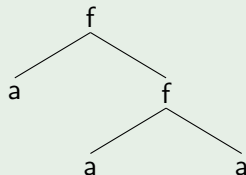
$c[t]$: the tree obtained by substituting $\$$ by t .

Example:

$$c = f(a, \$)$$



$$c[f(a, a)] = f(a, f(a, a))$$



An Algebraic Characterization of Rational Series

Contexts operate on tree series

Let $c \in C(\mathcal{F})$. Define $\dot{c} : \mathbb{R}\langle\langle T(\mathcal{F}) \rangle\rangle \rightarrow \mathbb{R}\langle\langle T(\mathcal{F}) \rangle\rangle$ by

$$\dot{c}r(t) = r(c[t]).$$

Example

$c = f(a, \$)$, $t = f(a, a)$, $\dot{c}r(t) = r(f(a, f(a, a)))$.

Let $r \in T(\mathcal{F})$, consider $W_r = [\{\dot{c}r \mid c \in C(\mathcal{F})\}] \subseteq \mathbb{R}\langle\langle T(\mathcal{F}) \rangle\rangle$
 the vector subspace of $\mathbb{R}\langle\langle T(\mathcal{F}) \rangle\rangle$ spanned by the series $\dot{c}r$.

Theorem: r is rational iff the dimension of W_r is finite.

The Canonical Linear Representation of Rational Series

$W_r = [\{\dot{c}r \mid c \in C(\mathcal{F})\}]$; W_r^* dual space of W_r

- No natural linear representation of r on W_r

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$$t \rightarrow \bar{t} \text{ s.t. } \bar{t}(\dot{c}r) = r(c[t])$$

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$$t \rightarrow \bar{t} \text{ s.t. } \bar{t}(\dot{c}r) = r(c[t])$$

- $\{\bar{t} \mid t \in T(\mathcal{F})\}$ spans W_r^*
- the **canonical linear representation** of r :
(W_r^*, μ, λ) where $\mu(t) = \bar{t}$ and $\lambda = r$ ($W_r^{**} = W_r$)

Building the Canonical Linear Representation

$$\mathcal{F} = \{f(,), a\}, \tau(q) = 1, p_\alpha : q \xrightarrow{\alpha} a, q \xrightarrow{1-\alpha} f(q, q)$$

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Let $p = 2p_{2/3} - p_{3/4} : \sum_t p(t) = 1$ and $\forall t, p(t) \geq 0$.

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$$p(a) = \frac{7}{12}, p(f(a, a)) = \frac{269}{1728}, p(f(a, f(a, a))) = p(f(f(a, a), a)) = \frac{9823}{248832}, \dots$$

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Oracle: Is $\bar{a} = 0$? i.e. for every context c , $p(c[a]) = 0$?

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Answer: **NO**, consider $c = \$$.

Let $B = \{\bar{a}\}$.

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i.e. $\exists \alpha$, for every context c , $p(c[f(a, a)]) = \alpha p(c[a])$?

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Answer: **NO**, consider $c_1 = \$$ and $c_2 = f(a, \$)$.

Let $B = \{\bar{a}, \overline{f(a, a)}\}$.

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Oracle: Is $\overline{f(a, f(a, a))}$ colinear to $\bar{a}, \overline{f(a, a)}$?

Answer: YES,

$$\overline{f(a, f(a, a))} = \frac{-54}{2^4 \times 3^4} \bar{a} + \frac{59}{2^4 \times 3^2} \overline{f(a, a)}.$$

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Answer: YES,

$$\overline{f(f(a, a), f(a, a))} = \frac{-3186}{2^8 \times 3^6} \bar{a} + \frac{2617}{2^8 \times 3^4} \overline{f(a, a)}.$$

Let $B = \{\bar{a}, \overline{f(a, a)}\}$.

Building the Canonical Linear Representation

$$p = 2p_{2/3} - p_{3/4}$$

$$B = \{\bar{a}, \overline{f(a, a)}\}.$$

$$\mu(a) = \bar{a}$$

$$\mu(f)(\bar{a}, \bar{a}) = \overline{f(a, a)}$$

$$\mu(f)(\bar{a}, \overline{f(a, a)}) = \frac{-54}{2^4 \times 3^4} \bar{a} + \frac{59}{2^4 \times 3^2} \overline{f(a, a)}$$

$$\mu(f)(\overline{f(a, a)}, \bar{a}) = \frac{-54}{2^4 \times 3^4} \bar{a} + \frac{59}{2^4 \times 3^2} \overline{f(a, a)}$$

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$$\lambda(\bar{a}) = p(a) = \frac{7}{12}; \lambda(\overline{f(a, a)}) = p(f(a, a)) = \frac{269}{1728}$$

Algorithm DEES; Independence Test

S a finite sample i.i.d. from p ; B current basis; s vector candidate

$$\forall \alpha_t \in \mathbb{R}, \bar{s} \neq \sum_{t \in B} \alpha_t \bar{t}$$

\simeq

$$\bigwedge_{c: \exists t \ c[t] \in S} \left\{ |p_S(c[s]) - \sum_{t \in B} \alpha_t p_S(c[t])| \leq \epsilon \right\} \text{ has no solution.}$$

Take $\epsilon = |S|^{-\gamma}$ where $\gamma \in]0, 1/2[$ (VC bounds).

Properties of DEES

Theorem [F. Denis and A. Habrard, ALT'07]

DEES identifies the correct basis in the limit with probability one and the parameters converge to the correct ones in $O(|S|^{-1/2})$.

But ...

- In the model output, the states may not define stochastic languages.
- The parameters are not normalized.
- Before convergence, the model output may not define a stochastic language.

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The Normalization of the Model

$$q \rightarrow q_0, 7/12 + q_1, 269/1728$$

$$q_0 \rightarrow a, 1 + f(q_0, q_1), \frac{-54}{2^4 3^4} + f(q_1, q_0), \frac{-54}{2^4 3^4} + f(q_1, q_1), \frac{-3186}{2^8 3^6}$$

$$q_1 \rightarrow f(q_0, q_0), 1 + f(q_0, q_1), \frac{59}{2^4 3^2} + f(q_1, q_0), \frac{59}{2^4 3^2} + f(q_1, q_1), \frac{2617}{2^8 3^4}$$

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$$q_1 \rightarrow f(q_0, q_0), 1 + f(q_0, q_1), \frac{59}{2^4 3^2} + f(q_1, q_0), \frac{59}{2^4 3^2} + f(q_1, q_1), \frac{2617}{2^8 3^4}$$

Theorem

For any rational stochastic language, there exists a normalized representation with a basis chosen to ensure that:

- Each state defines a stochastic language.
- The weights of the transitions are normalized.

After Renormalization

- \forall state lhs: Sum of the transition weights is one.
- \forall pair (state-lhs,symbol): Sum of the transition weights ≥ 0 .

$$\begin{aligned}
 q &\rightarrow q_0, 1 \\
 q_0 &\rightarrow a, \frac{7}{12} + f(q_0, q_0), \frac{-269}{50} + f(q_0, q_1), \frac{259}{50} + f(q_1, q_0), \frac{259}{50}, \\
 &\quad + f(q_1, q_1), \frac{-1369}{300} \\
 q_1 &\rightarrow a, \frac{269}{444} + f(q_0, q_0), \frac{-3024}{925} + f(q_0, q_1), \frac{2664}{925} + f(q_1, q_0), \frac{2664}{925} \\
 &\quad + f(q_1, q_1), \frac{-23273}{11100}
 \end{aligned}$$

- Efficient propagative method for computing the normalization.
- Still negative weights \rightarrow specific generation algorithm.

Notion of Strong Consistency

Rational Stochastic Tree Language Strongly Consistent

$$\text{Bounded average tree size: } \sum_t \rho(t)|t| < \infty$$

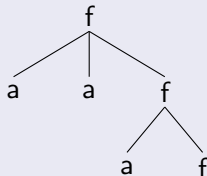
Theorem

For a strongly consistent RSTL, the spectral radius of the "expectation matrix" A taken from the normalized representation is strictly less than 1 ($\rho(A) < 1$).

Errata: Some hypotheses are missing in Proposition 1 see <http://hal.archives-ouvertes.fr/hal-00293511/en> (the series $\sum_{t \in T(\mathcal{F})} p_i(t)$ and $\sum_{t \in T(\mathcal{F})} p_i(t)|t|$ have to be absolutely convergent)

Adapting the Framework to Unranked Trees

Unranked Trees

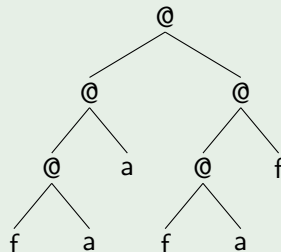


- Unranked tree series

\Leftrightarrow
Bijection

Ranked Trees

$$\mathcal{F}_0 = \{a, f\} \quad \mathcal{F}_2 = \{\textcircled{\text{a}}\}$$



- Ranked tree series

\Leftrightarrow
Equivalence

All the inference results apply: Convert the data and use DEES

Conclusion: Learning RSTL from *i.i.d.* samples

- DEES may output irrelevant representations.
- Our contributions:
 - Existence and construction of a normalized representation.
 - Algorithm for generating trees from the distribution.
 - Strong consistency.
 - Application to unranked trees.
- \Rightarrow When the models do not define stochastic languages, a distribution can be extracted and controlled if $\rho(A) < 1$.
- \Rightarrow A prototype software is being developed (Piccata).